

COMP 681 Formal Methods Spring 2008 Recitation 4—Solutions

1. Consider the following conditional. Note that the conditions of the conditionals within the {...} are mutually exclusive.

```

if P and Q then {
  if R then S1
  else S2
}
else if T then {
  if Q then S1
  else S2
}
else S2

```

- a. In the language of propositional logic, write the conditions under which each of the statements S_1 and S_3 is executed. Simplify, at least move all \neg 's inward so they apply only to primes and factor out common terms.

Answer

S_1

$$\begin{aligned}
 & P \wedge Q \wedge R \vee \neg(P \wedge Q) \wedge T \wedge Q \\
 \langle \equiv \rangle & P \wedge Q \wedge R \vee (\neg P \vee \neg Q) \wedge T \wedge Q && \text{De Morgan's 1st Law} \\
 \langle \equiv \rangle & P \wedge Q \wedge R \vee \neg P \wedge T \wedge Q \vee \neg Q \wedge T \wedge Q && \text{Distributive Law } (\wedge \text{ over } \vee) \\
 \langle \equiv \rangle & P \wedge Q \wedge R \vee \neg P \wedge T \wedge Q \vee T \wedge Q \wedge \neg Q && \text{Commutativity of } \wedge \\
 \langle \equiv \rangle & P \wedge Q \wedge R \vee \neg P \wedge T \wedge Q \vee T \wedge \text{false} && \text{Law of Contradiction} \\
 \langle \equiv \rangle & P \wedge Q \wedge R \vee \neg P \wedge T \wedge Q \vee \text{false} && \text{Simplification (3.20)} \\
 \langle \equiv \rangle & P \wedge Q \wedge R \vee \neg P \wedge T \wedge Q && \text{Simplification (3.21)} \\
 \langle \equiv \rangle & Q \wedge P \wedge R \vee Q \wedge \neg P \wedge T && \text{Commutativity of } \wedge (2\times) \\
 \langle \equiv \rangle & Q \wedge (P \wedge R \vee \neg P \wedge T) && \text{Distributive Law } (\wedge \text{ over } \vee)
 \end{aligned}$$

S_2

$$\begin{aligned}
 & P \wedge Q \wedge \neg R \vee \neg(P \wedge Q) \wedge T \wedge \neg Q \vee \neg(P \wedge Q) \wedge \neg T \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee (\neg P \vee \neg Q) \wedge T \wedge \neg Q \vee \neg(P \wedge Q) \wedge \neg T && \text{De Morgan's 1st Law} \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge T \wedge \neg Q \vee \neg Q \wedge T \wedge \neg Q \vee \neg(P \wedge Q) \wedge \neg T && \text{Distr. } (\wedge \text{ over } \vee) \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge T \wedge \neg Q \vee T \wedge \neg Q \wedge \neg Q \vee \neg(P \wedge Q) \wedge \neg T && \text{Comm. of } \wedge \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge T \wedge \neg Q \vee T \wedge \neg Q \vee \neg(P \wedge Q) \wedge \neg T && \text{Idempotence of } \wedge \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee T \wedge \neg Q \vee \neg(P \wedge Q) \wedge \neg T && \text{Absorption (3.22)} \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee T \wedge \neg Q \vee (\neg P \vee \neg Q) \wedge \neg T && \text{De Morgan's 1st Law} \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee T \wedge \neg Q \vee \neg P \wedge \neg T \vee \neg Q \wedge \neg T && \text{Distr. } (\wedge \text{ over } \vee) \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \vee T \wedge \neg Q \vee \neg Q \wedge \neg T && \text{Commutativity of } \vee \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \vee \neg Q \wedge T \vee \neg Q \wedge \neg T && \text{Commutativity of } \wedge \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \vee \neg Q \wedge (T \vee \neg T) && \text{Distr. } (\wedge \text{ over } \vee) \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \vee \neg Q \wedge \text{true} && \text{Excluded Middle} \\
 \langle \equiv \rangle & P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \vee \neg Q && \text{Simplification (3.18)}
 \end{aligned}$$

b. Rewrite the conditional in the form

```

if _ then {
  if _ then {
    if _ then ___
    else ___
  }
  else {
    if _ then ___
    else ___
  }
}
else ___

```

Answer

First modify the expression for S_2 as follows:

$P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \vee \neg Q$	The result above
$\langle \equiv \rangle P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \wedge true \vee \neg Q$	Simplification (3.18)
$\langle \equiv \rangle P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \wedge (Q \vee \neg Q) \vee \neg Q$	Law of Excluded Middle
$\langle \equiv \rangle P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \wedge Q \vee \neg P \wedge \neg T \wedge \neg Q \vee \neg Q$	Distr. (\wedge over \vee)
$\langle \equiv \rangle P \wedge Q \wedge \neg R \vee \neg P \wedge \neg T \wedge Q \vee \neg Q$	Idempotence (3.22)

Using the final expression above for S_1 and the expression for S_2 derived here, we can read off the conditions.

```

if Q then {
  if P then {
    if R then S1
    else S2
  }
  else {
    if T then S1
    else S2
  }
}
else S2

```

2. A group of experts from different departments decides that procedure S_1 is done when P is true, that S_2 is done if R is false, and that S_3 is done if both Q and R hold.

a. Write a GCL alternation statement that captures these requirements. (The conditions won't be mutually exclusive at this point.)

Answer

```

if    P  →  S1
□    ¬ R  →  S2
□    Q ∧ R →  S3
fi

```

- b. You point out to the experts that their conditions are not mutually exclusive. They decide among themselves that S_1 should take priority over the other two statements and that S_3 should take priority over S_2 . Rewrite the conditions so that they are mutually exclusive and respect these priorities.

Answer

Before simplification, we have

```

if      P                               →   S1
□  ¬ P ∧ ¬ (Q ∧ R) ∧ ¬ R             →   S2
□  ¬ P ∧ Q ∧ R                         →   S3
fi

```

We simplify the guard for S_2 as follows.

```

¬ P ∧ ¬ (Q ∧ R) ∧ ¬ R
⟨≡⟩ ¬ P ∧ (¬ Q ∨ ¬ R) ∧ ¬ R    De Morgan's 1st Law
⟨≡⟩ ¬ P ∧ ¬ R                  Absorption (3.23)

```

So the above simplifies to

```

if      P                               →   S1
□  ¬ P ∧ ¬ R                             →   S2
□  ¬ P ∧ Q ∧ R                             →   S3
fi

```

- c. For the cases not covered in **b**, add another alternative whose guard covers the remaining cases yet is disjoint from each of the other three guards. The command associated with this guard should be `skip`.

Answer

We form the conjunction of the negations of the guards and simplify:

$$\begin{aligned}
 & \neg P \wedge \neg(\neg P \wedge \neg R) \wedge \neg(\neg P \wedge Q \wedge R) \\
 \langle \equiv \rangle & \neg P \wedge (\neg\neg P \vee \neg\neg R) \wedge (\neg\neg P \vee \neg Q \vee \neg R) && \text{De Morgan's 1st Law (2}\times\text{)} \\
 \langle \equiv \rangle & \neg P \wedge (P \vee R) \wedge (P \vee \neg Q \vee \neg R) && \text{Law of Negation (3}\times\text{)} \\
 \langle \equiv \rangle & (\neg P \wedge P \vee \neg P \wedge R) \wedge (P \vee \neg Q \vee \neg R) && \text{Distributive Law } (\wedge \text{ over } \vee) \\
 \langle \equiv \rangle & (\text{false} \vee \neg P \wedge R) \wedge (P \vee \neg Q \vee \neg R) && \text{Law of Contradiction} \\
 \langle \equiv \rangle & \neg P \wedge R \wedge (P \vee \neg Q \vee \neg R) && \text{Simplification (3.21)} \\
 \langle \equiv \rangle & \neg P \wedge R \wedge P \vee \neg P \wedge R \wedge \neg Q \vee \neg P \wedge R \wedge \neg R && \text{Distributive Law } (\wedge \text{ over } \vee) \\
 \langle \equiv \rangle & \neg P \wedge R \wedge P \vee \neg P \wedge R \wedge \neg Q \vee \neg P \wedge \text{false} && \text{Law of Contradiction} \\
 \langle \equiv \rangle & \neg P \wedge R \wedge P \vee \neg P \wedge R \wedge \neg Q \vee \text{false} && \text{Simplification (3.20)} \\
 \langle \equiv \rangle & \neg P \wedge R \wedge P \vee \neg P \wedge R \wedge \neg Q && \text{Simplification (3.21)} \\
 \langle \equiv \rangle & R \wedge P \wedge \neg P \vee \neg P \wedge R \wedge \neg Q && \text{Commutativity of } \wedge \\
 \langle \equiv \rangle & R \wedge \text{false} \vee \neg P \wedge R \wedge \neg Q && \text{Law of Contradiction} \\
 \langle \equiv \rangle & \text{false} \vee \neg P \wedge R \wedge \neg Q && \text{Simplification (3.20)} \\
 \langle \equiv \rangle & \neg P \wedge R \wedge \neg Q && \text{Simplification (3.21)} \\
 \langle \equiv \rangle & \neg P \wedge \neg Q \wedge R && \text{Law of Contradiction}
 \end{aligned}$$

So the final alternation statement is

```

if    P           →  S1
□  ¬ P ∧ ¬ R     →  S2
□  ¬ P ∧ Q ∧ R   →  S3
□  ¬ P ∧ ¬ Q ∧ R →  skip
fi

```

- d. Translate the alternative statement in **c** into the language used in section 3.3
Minimize the number of primes that must be evaluated.

Answer

```

if P then S1
else {
  if R then {
    if Q then S3
  }
  else S2
}

```